

COQ

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Background



- Based on the Calculus of Inductive Constructions
 - Machine-checked proofs for software and math

Defining logical operators [\[edit \]](#)

The calculus of constructions has very few basic operators: the only logical operator for forming propositions is \forall . However, this one operator is sufficient to define all the other logical operators:

$$\begin{aligned} A \Rightarrow B &\equiv \forall x : A. B && (x \notin B) \\ A \wedge B &\equiv \forall C : \mathbf{P}. (A \Rightarrow B \Rightarrow C) \Rightarrow C \\ A \vee B &\equiv \forall C : \mathbf{P}. (A \Rightarrow C) \Rightarrow (B \Rightarrow C) \Rightarrow C \\ \neg A &\equiv \forall C : \mathbf{P}. (A \Rightarrow C) \\ \exists x : A. B &\equiv \forall C : \mathbf{P}. (\forall x : A. (B \Rightarrow C)) \Rightarrow C \end{aligned}$$

“If A then B” = For any assumption x of type A, B must follow

Background



- Runs on Gallina
 - Functional and strongly typed language
 - [Specification](#) language and programming language

Verifying $0 + n = n$

1. Prove that for all natural numbers adding zero to n results in n
2. Start proof block
3. Assume that n is an arbitrary natural number
4. Simplify the expression
5. We have deduced that both sides of the equation are equal, the proof is verified
6. End proof block

```
1 Theorem zero_plus_n : forall n:nat, 0 + n = n.  
2 Proof.  
3   intros n.  
4   simpl.  
5   reflexivity.  
6 Qed.
```

Failing proof

3. Define recursive (fixpoint) function (ack) to take two natural numbers as the input ($m\ n : \text{nat}$)
4. Pattern match on a pair of natural numbers (if statements)
5. If $m = 0$, return $n + 1$ ($_$ is wildcard)
6. If $m > 0$ and $n = 0$, recursively call ack
7. If $m > 0$ and $n > 0$, call ack with an arg to call the function again (twice)

Nested recursion will fail and function will never terminate (not structural recursion)

Require Import Arith.

```
Fixpoint ack (m n : nat) : nat :=  
  match m, n with  
  | 0, _ => n + 1  
  | S m', 0 => ack m' 1  
  | S m', S n' => ack m' (ack m' n')  
end.
```

Modified: Coq Development Team. “The Ackermann Function.” *The Coq Proof Assistant, Reference Manual*, v 8.19 (Inria, 2024), § Extraction, “The Ackermann Function.” Accessed April 22, 2025.
<https://rocq-prover.org/doc/v8.19/refman/proofs/automatic-tactics/auto.html?highlight=ack>

Why use it?

TLA+:

- Formal specifications of systems
- Model concurrent system behavior

Think: temporal model checking

THEOREM PlusZero == $\forall n \in \text{Nat} : n + 0 = n$

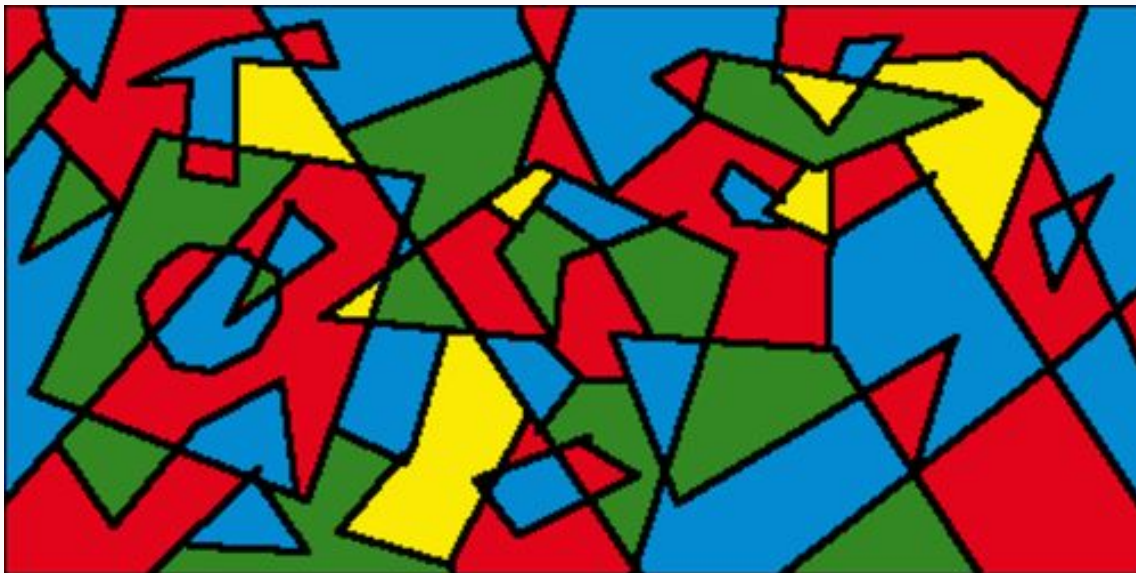
Coq:

- Formal proofs about programs and Math
- Verifying code logic and correctness

Think: machine-verified using type theory

Lemma plus_0_r (n : nat) : n + 0 = n. Proof. now
rewrite Nat.add_0_r. Qed.

Why use it? - ext



- Four Color Theorem
 - Proven in 2008 with Coq that any partitioned map can be colored with 4 colors and not touch any shape with the same color (minus corners)

Rating - Pros

- Expressiveness: proof automation + extraction to OCaml, Haskell (4/5)
- Well-definedness: Calculus of Inductive Constructions (5/5)
- Readability: can be intuitive to learn with those with a math/logic background (3.5/5)
- Reliability: used in CompCert C compiler (5/5)

If you value a lightweight, simple solution to logical verification, this is for you

Rating - Cons

- Very old formal method language, there are better, newer options
- More Math than Software Engineering proving

Very useful for logical conditions, not so much for temporal logic / concurrency

Additional Reading

- [Certified Programming with Dependent Types: A Pragmatic Introduction to the Coq Proof Assistant](#) - 2013 pdf by Adam Chlipala, MIT Press Direct
- [Kaiyu Yang, Jia Deng - Learning to Prove Theorems via Interacting with Proof Assistants](#) - CoqGym, 71K compiled human-written proofs to train a Coq proof assistant
- CoqPyt - Python framework for interacting with Coq
- CoqPilot - VSC Extension for LLM-generated Coq proofs